



ulm university universität  
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Programmierung und Ausführung  
von taskbasierten HPC Anwendungen

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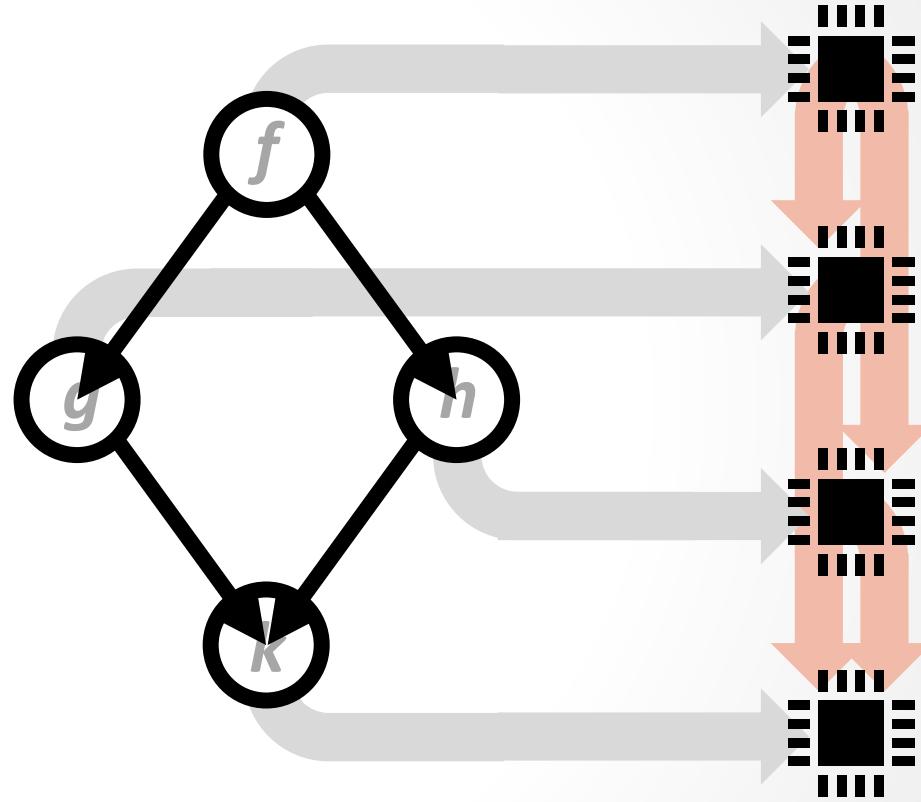
*Lutz Schubert*  
*[lutz.Schubert@uni-ulm.de](mailto:lutz.Schubert@uni-ulm.de)*

# Data Parallelism

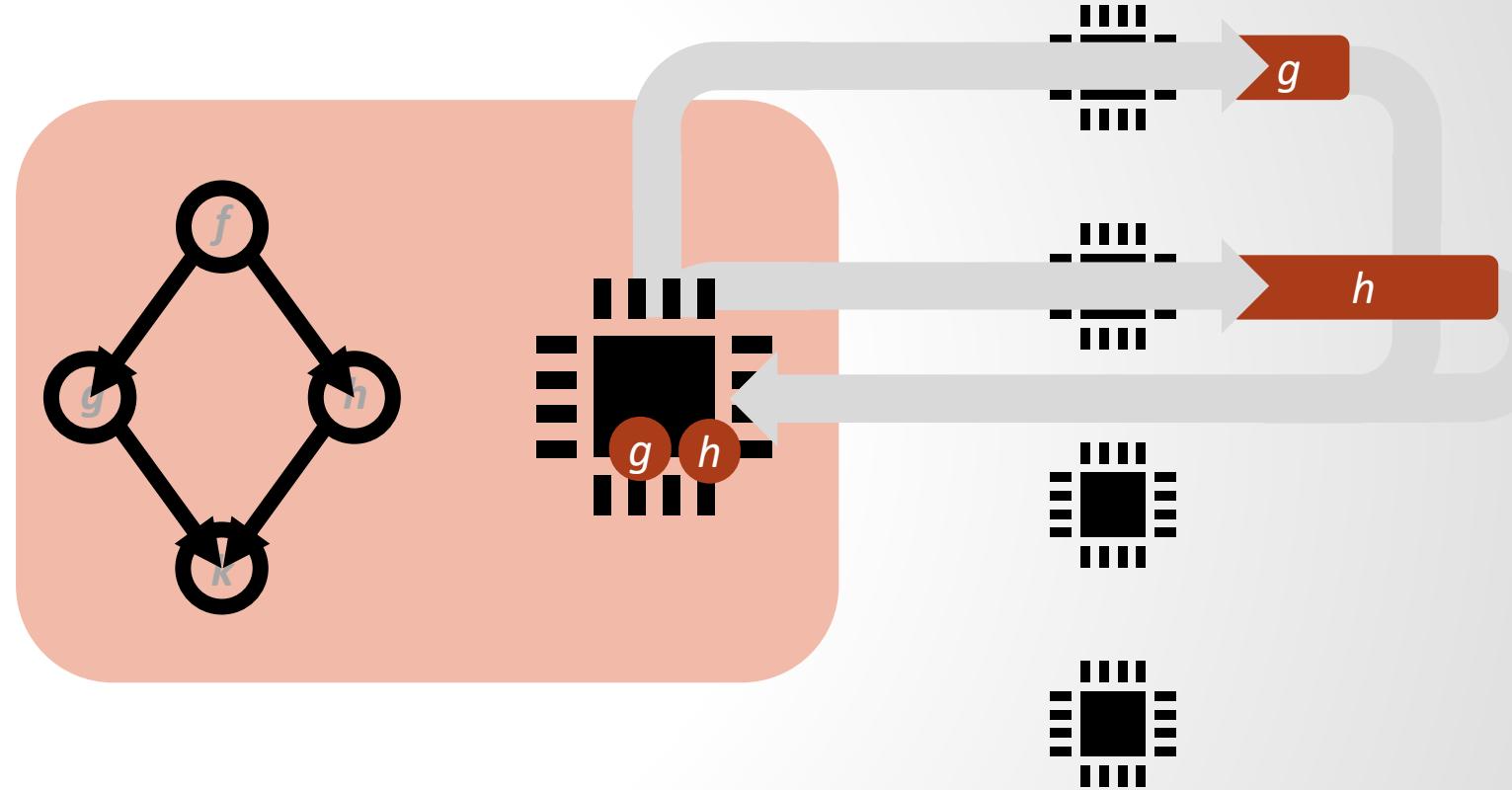
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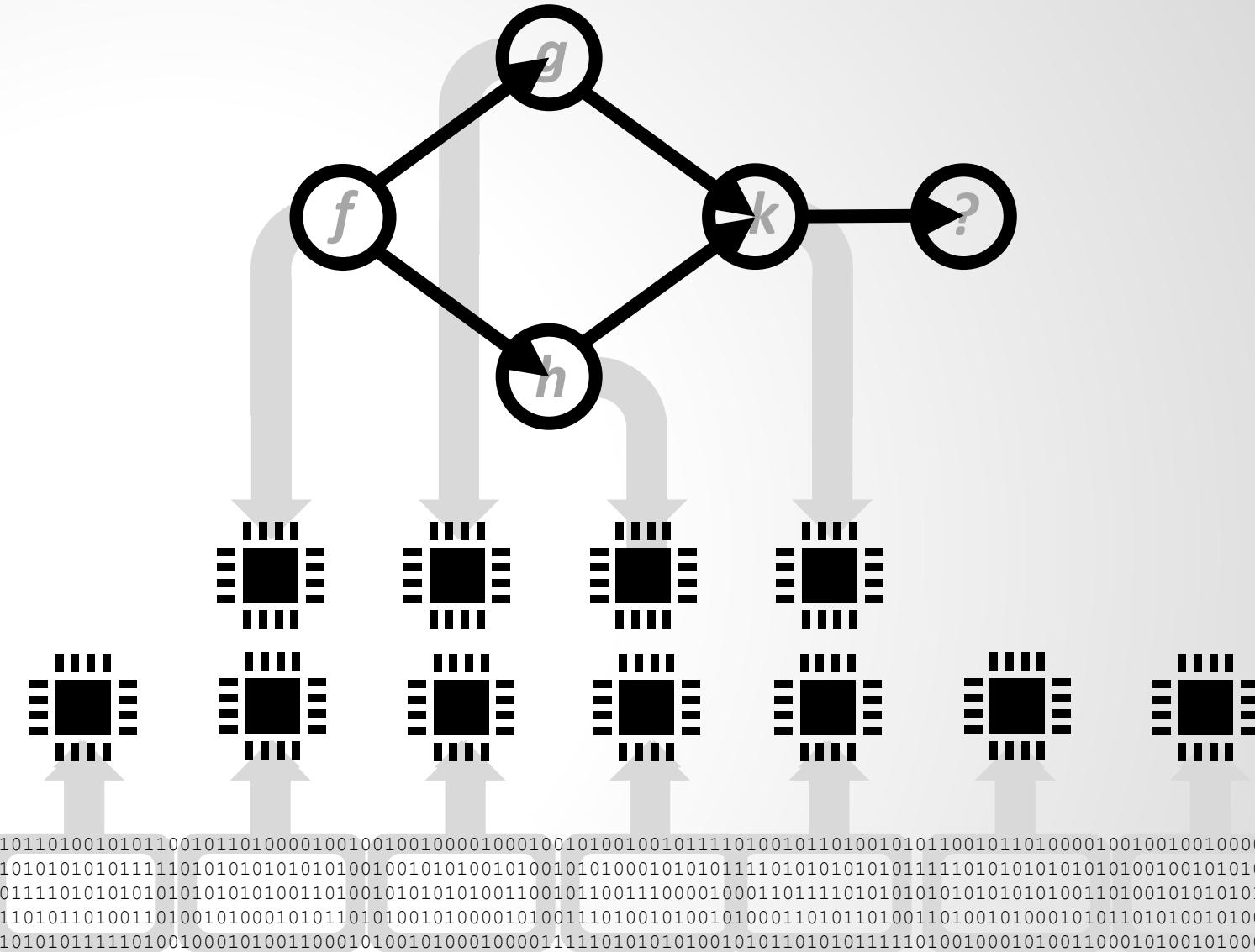
# Task-Based Parallelism



# Task- vs. Data-Based Parallelism Overhead



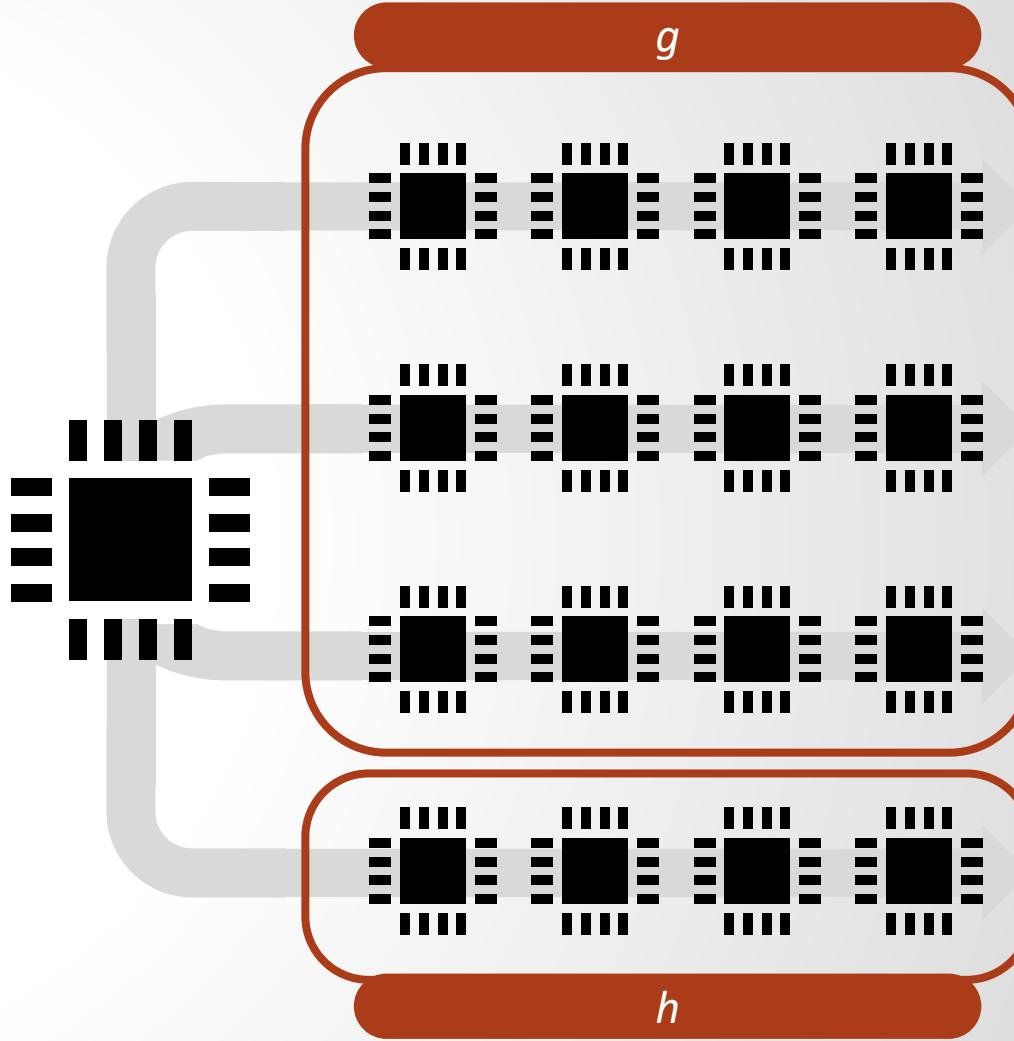
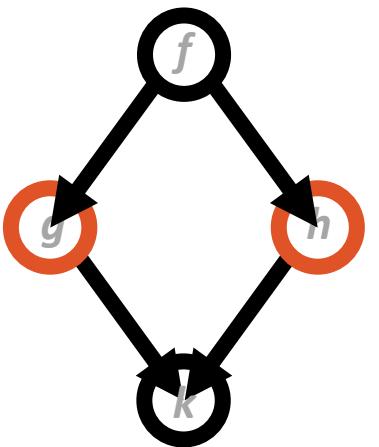
# Task- vs. Data-Based Parallelism Scale



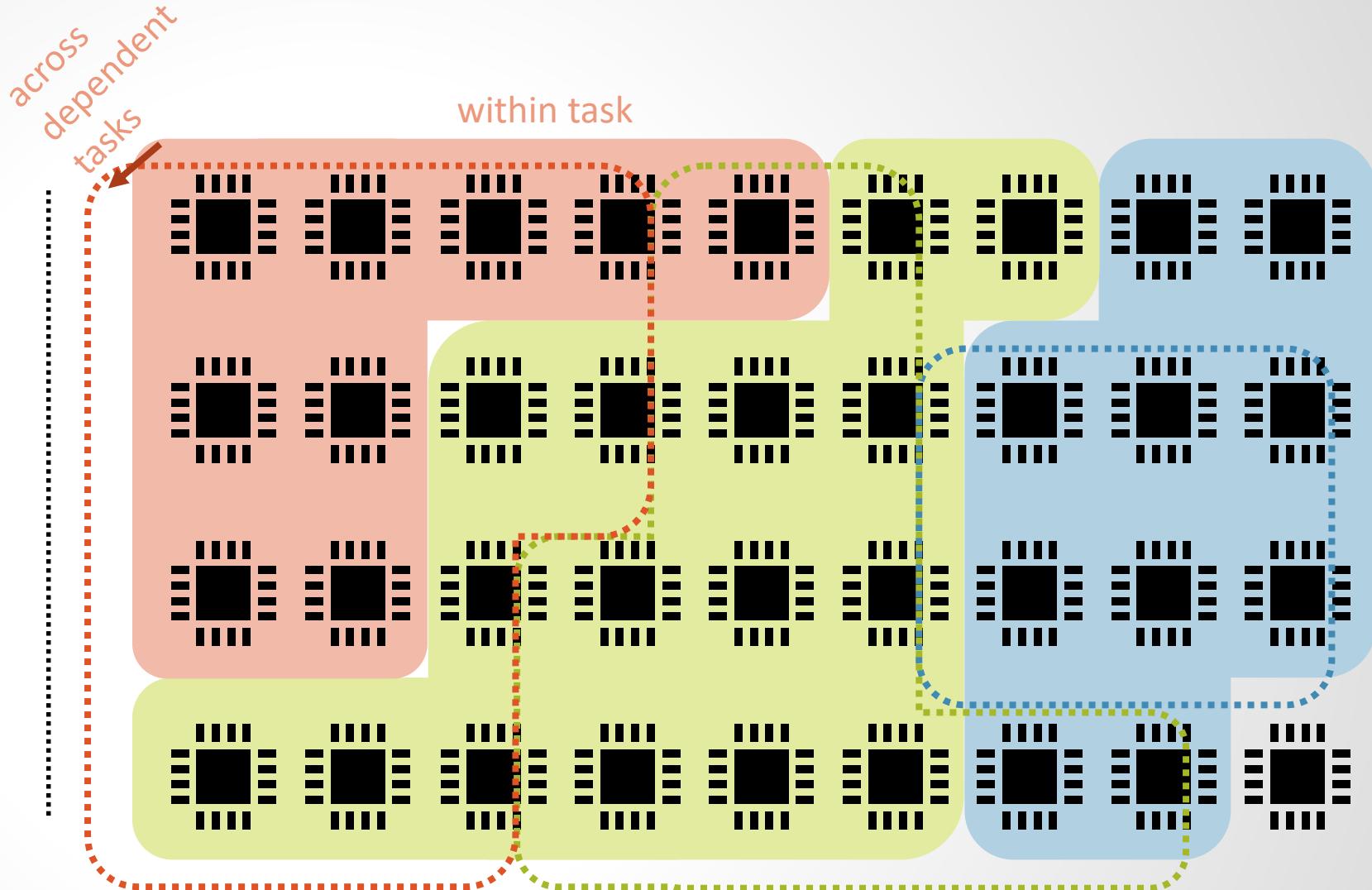
# Task- vs. Data- Based Parallelism

- All parallelisation generates overhead
- In general, overhead for task-based parallelism is higher than for data-based
  - Scheduling
  - Data gathering
  - Identification of next tasks
  - ...
- Most algorithms show a higher degree of data parallelism than task independency

# Mixed Parallelism



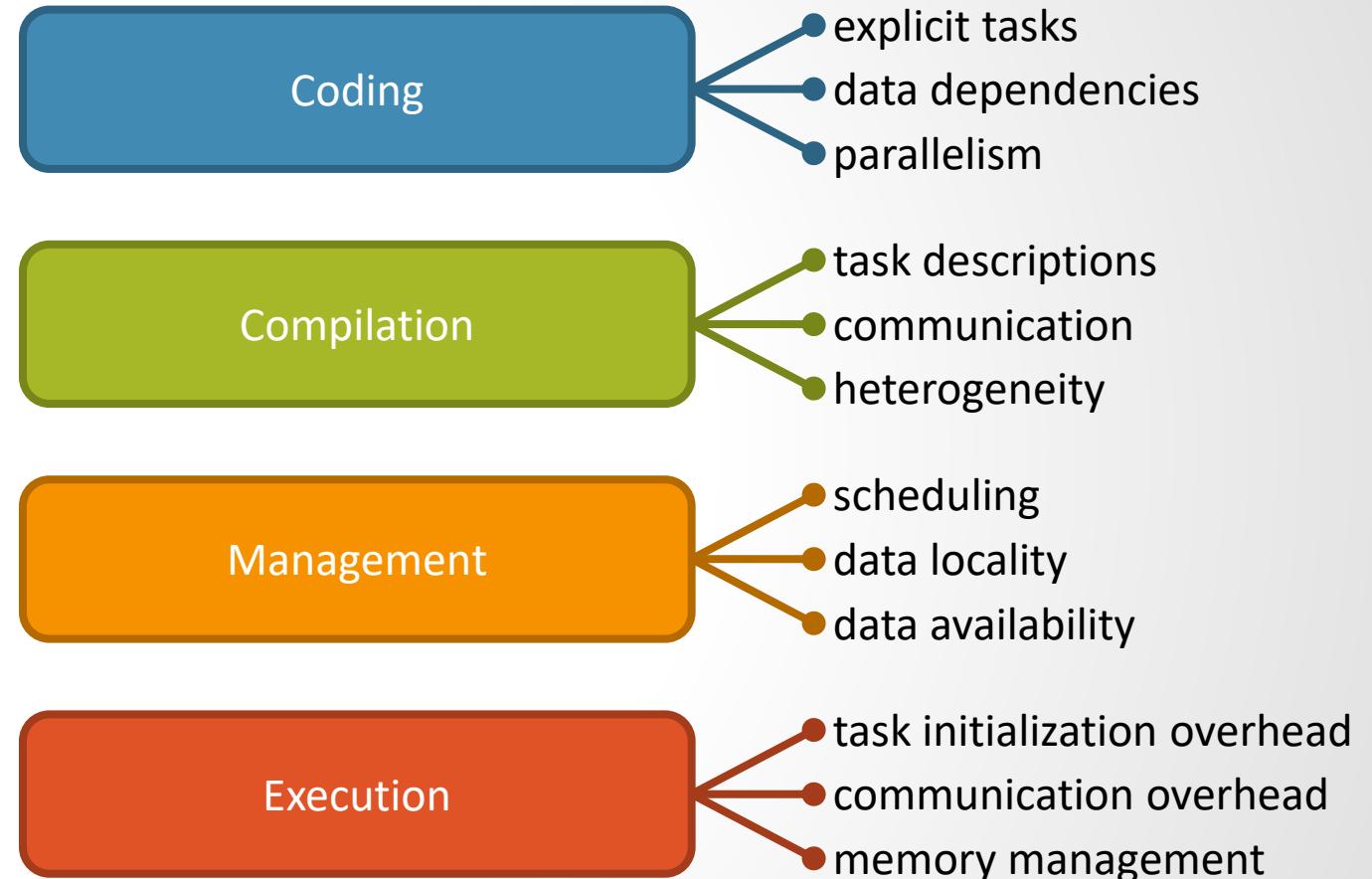
# Mixed Parallelism Maintaining Data Locality



# Mixed Parallelism Problems

- Complete restructuring of code is necessary
- Data dependencies need to be known
- Overhead must be minimal to allow for management of small tasks and threads

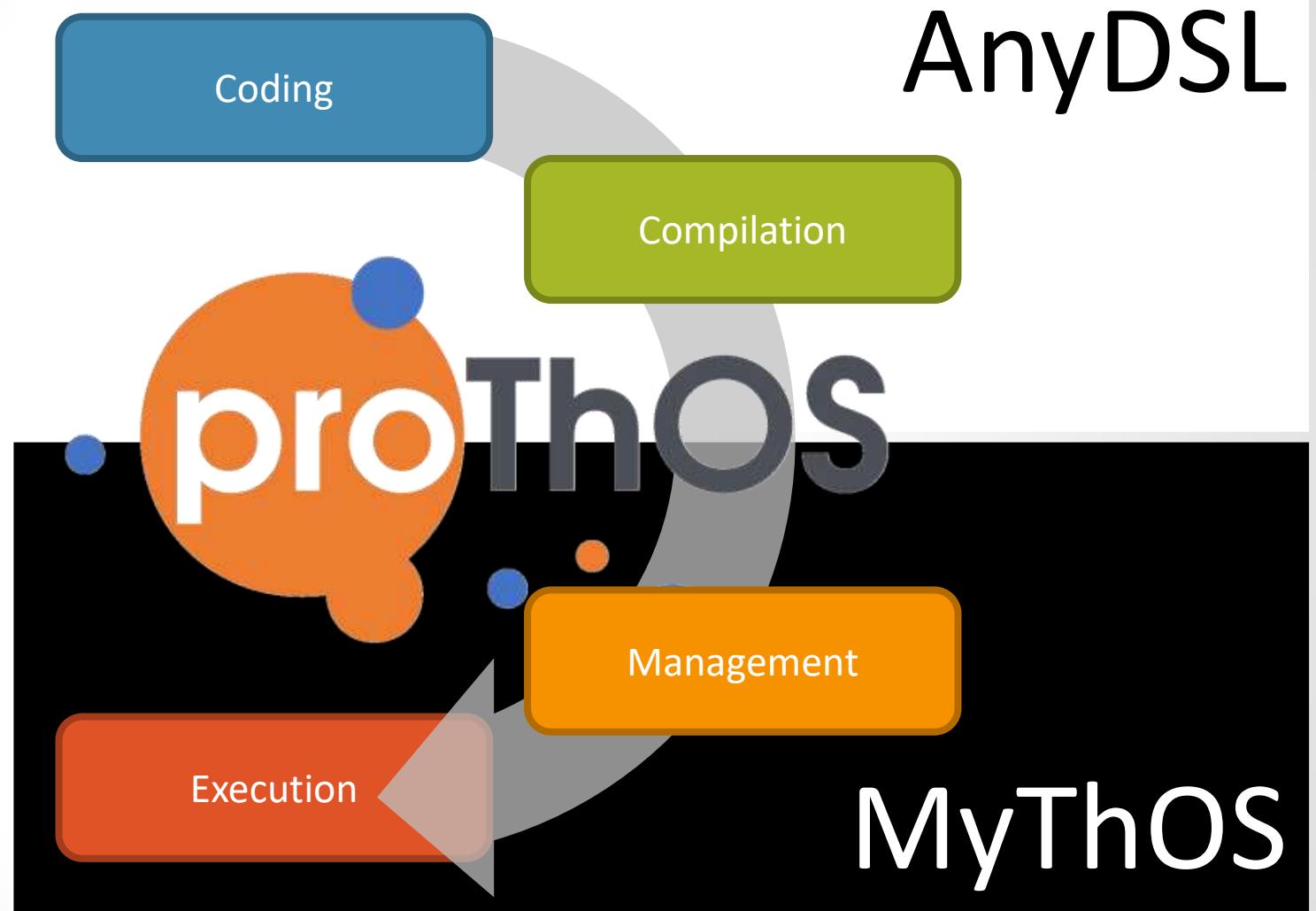
# Challenges for ProThOS



# General Scope



Two Major  
Building  
Blocks



# Two Major Building Blocks

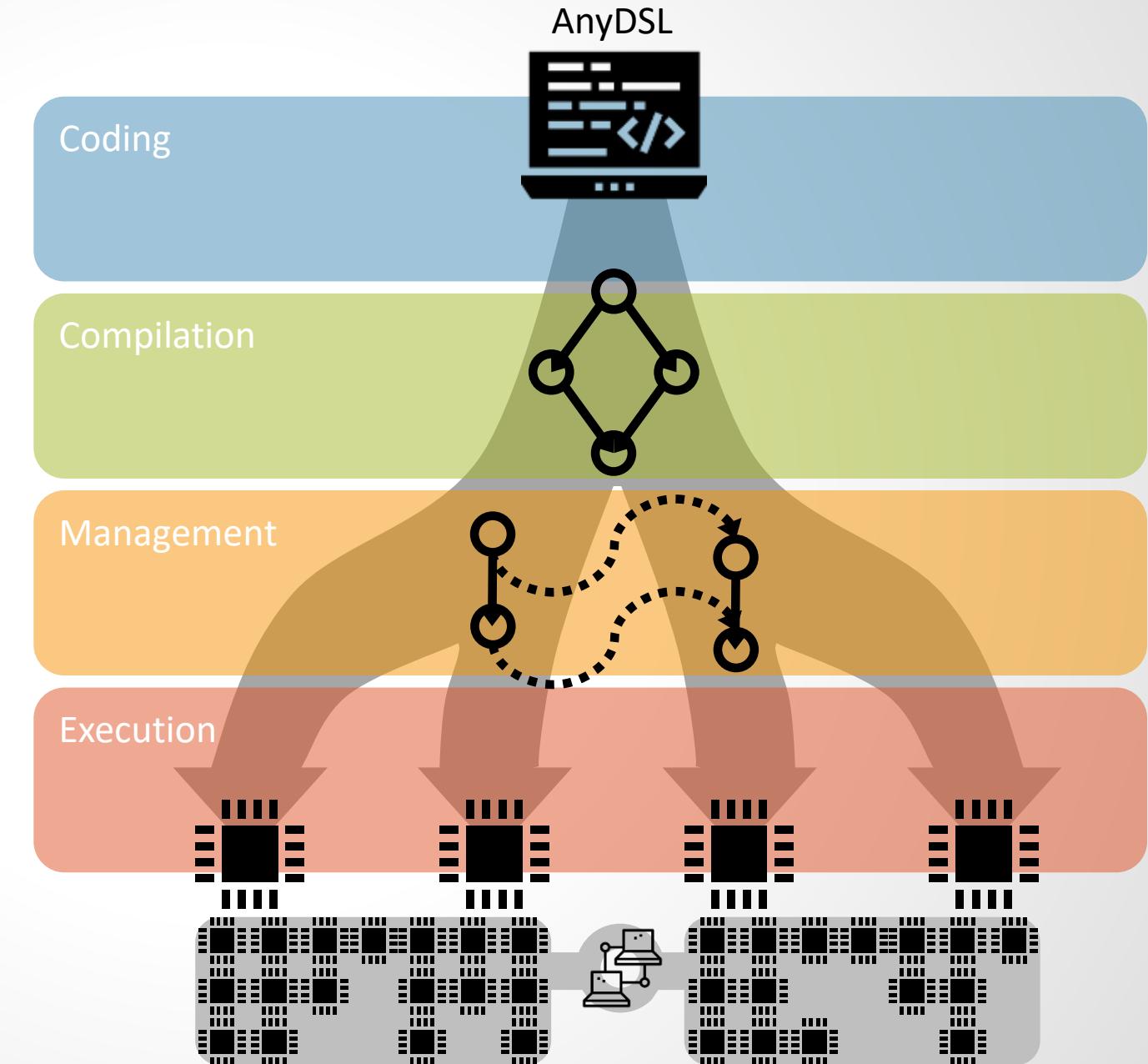
- Developed by DFKI & Uni Saarbrücken
- Allows to generate DSLs
- Highly efficient compilation
- Specifically developed for Vectorization

AnyDSL

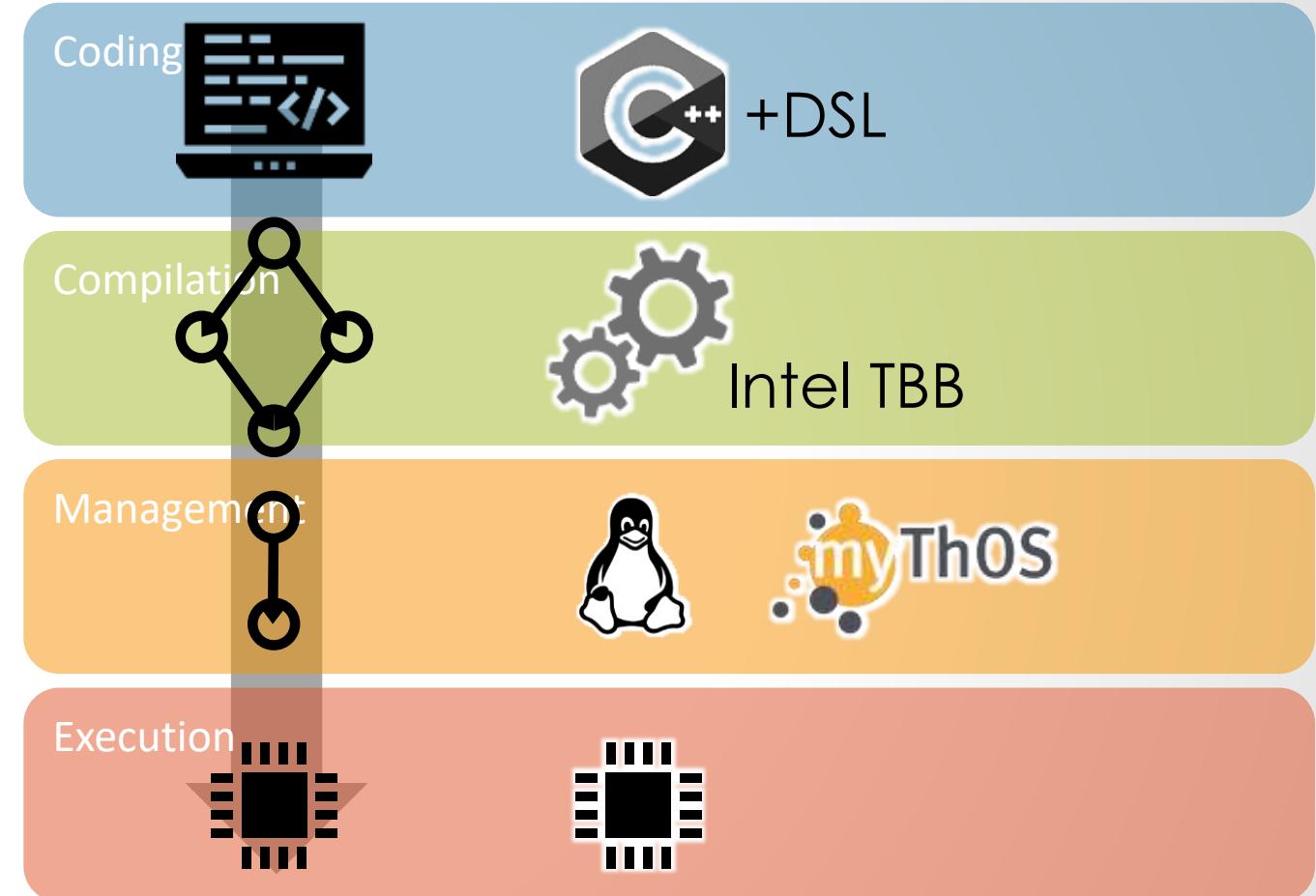
- Developed by BTU Cottbus and Uni Ulm
- Minimal, modular kernel
- Allows for easy extension with new functionalities
- Much better pThread management than Linux

MyThOS

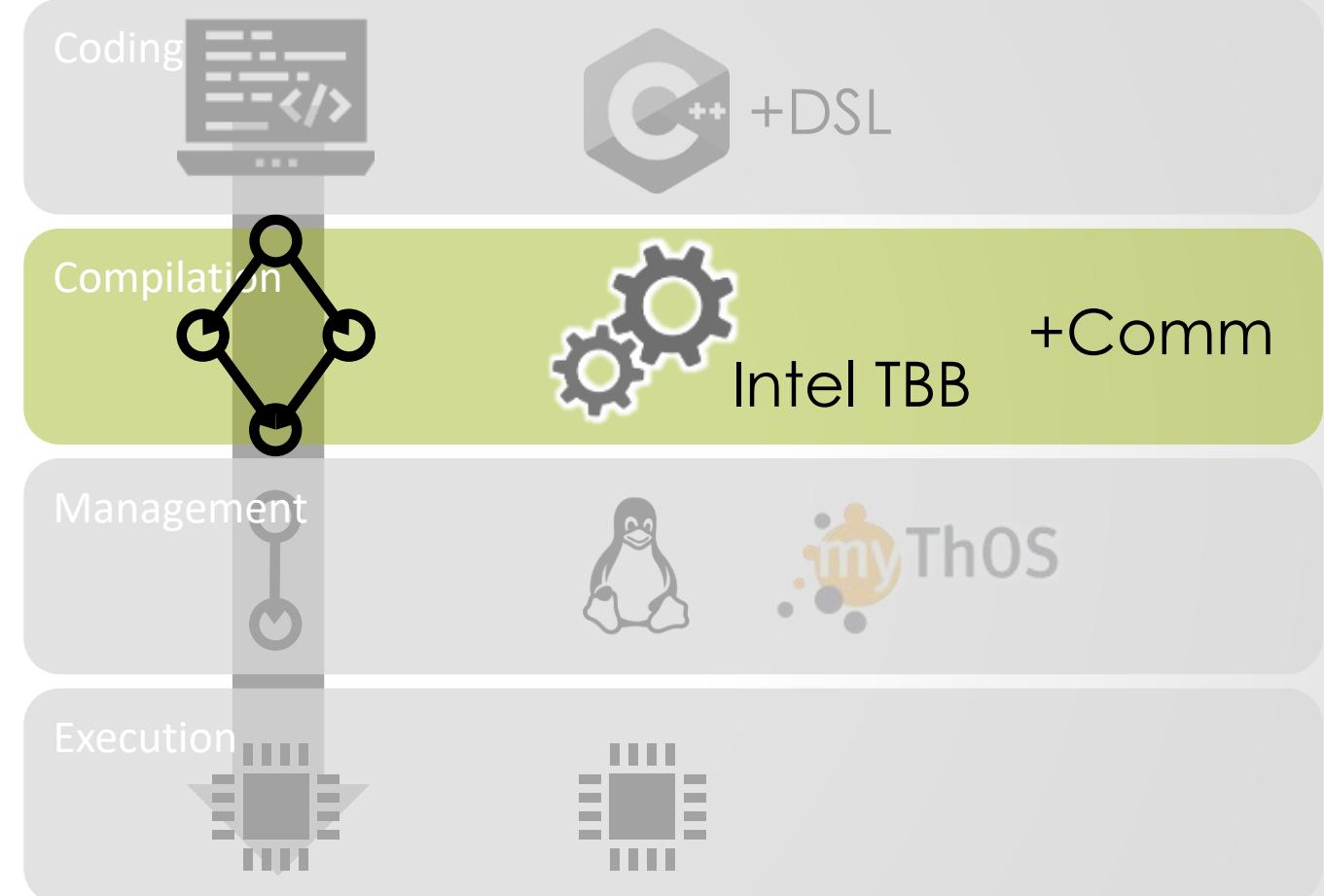
# ProThOS Lifecycle



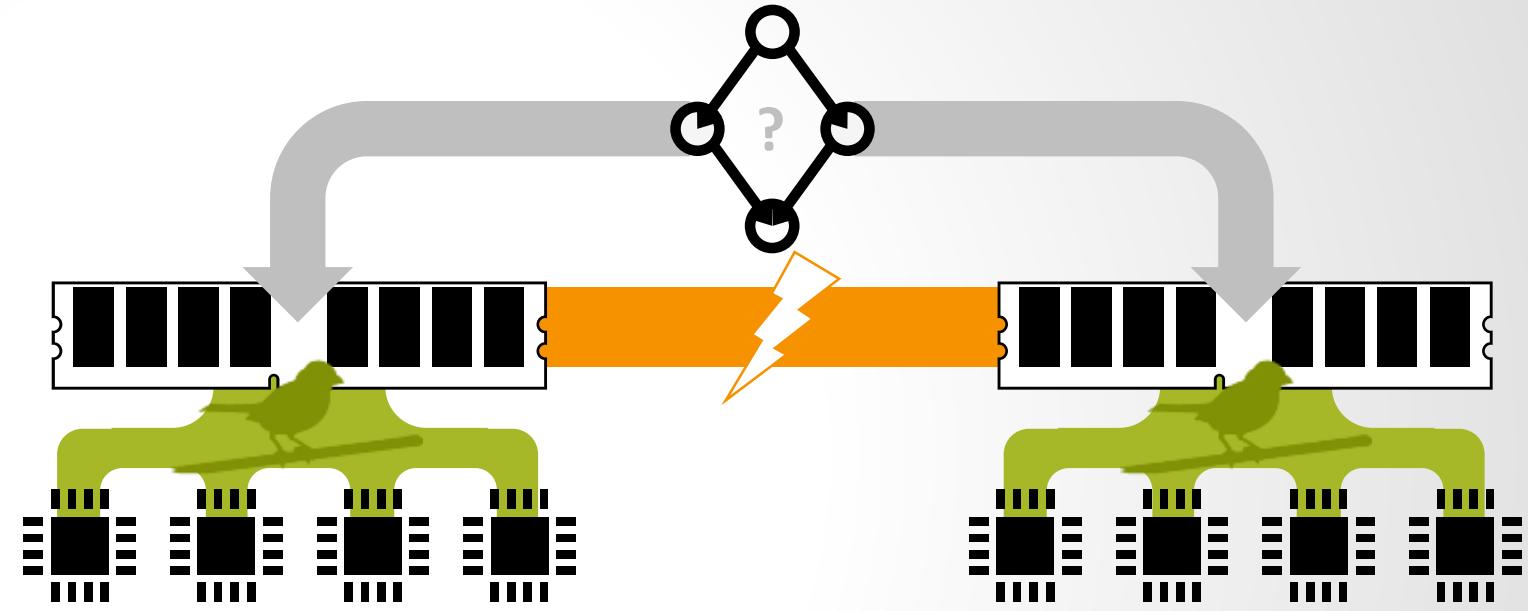
# Principle Technology Support



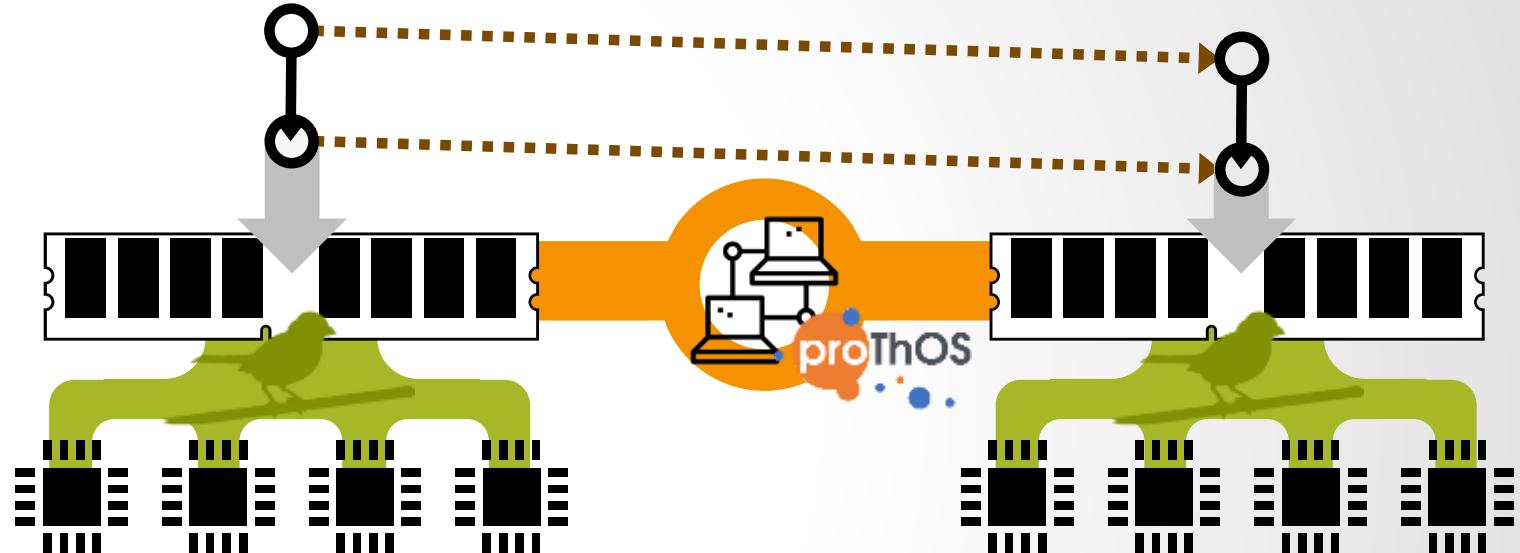
# Actual Technology & Status Intel TBB



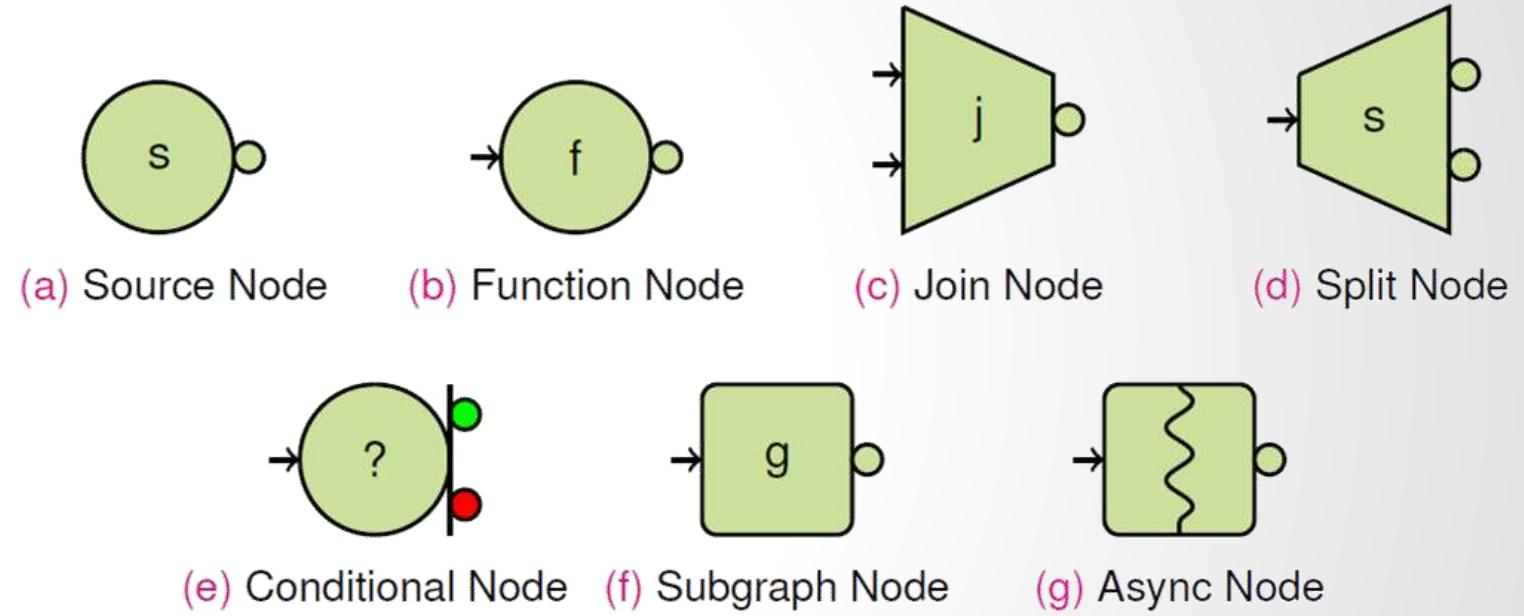
# Actual Technology & Status Task Description



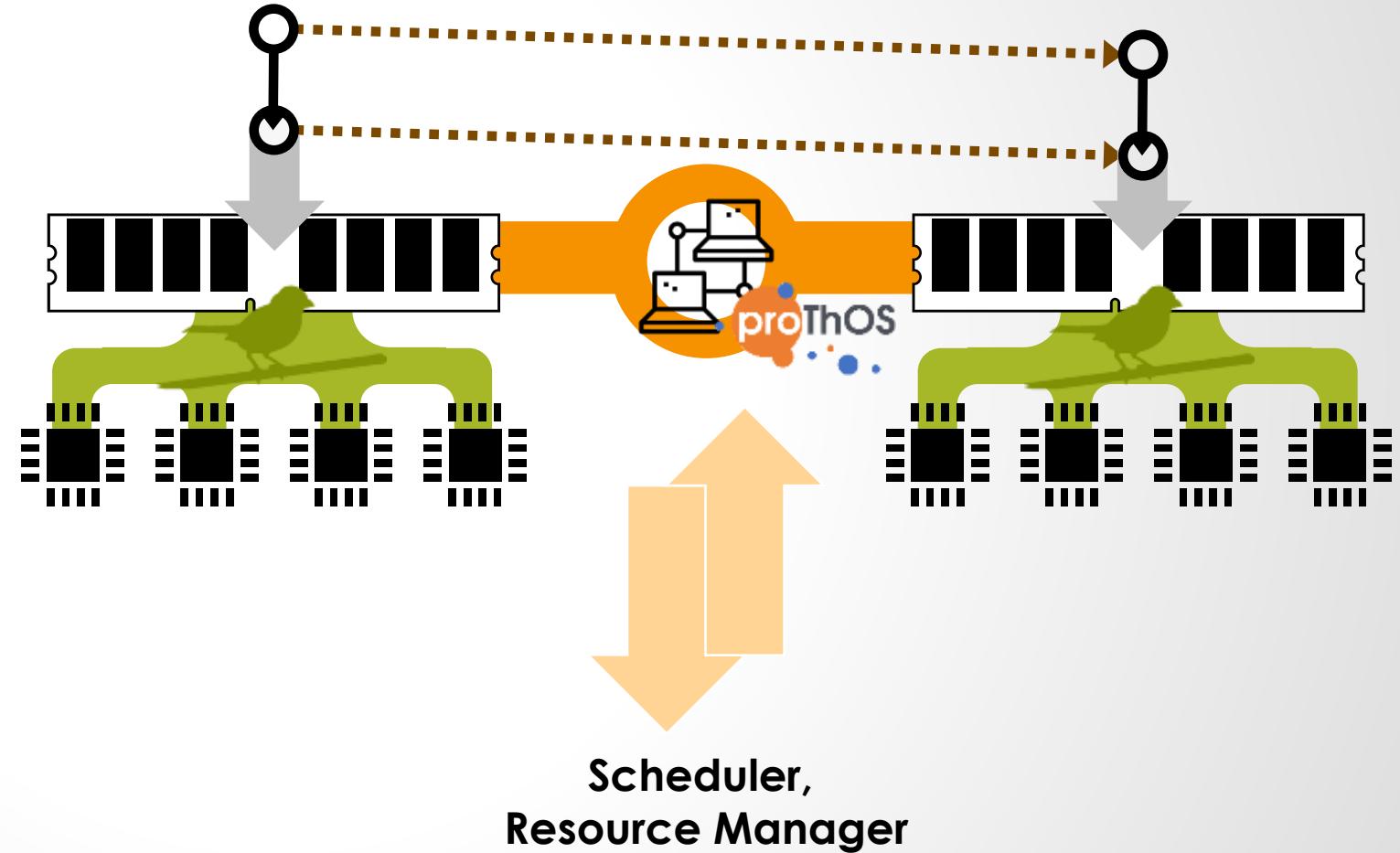
# Actual Technology & Status Task Description



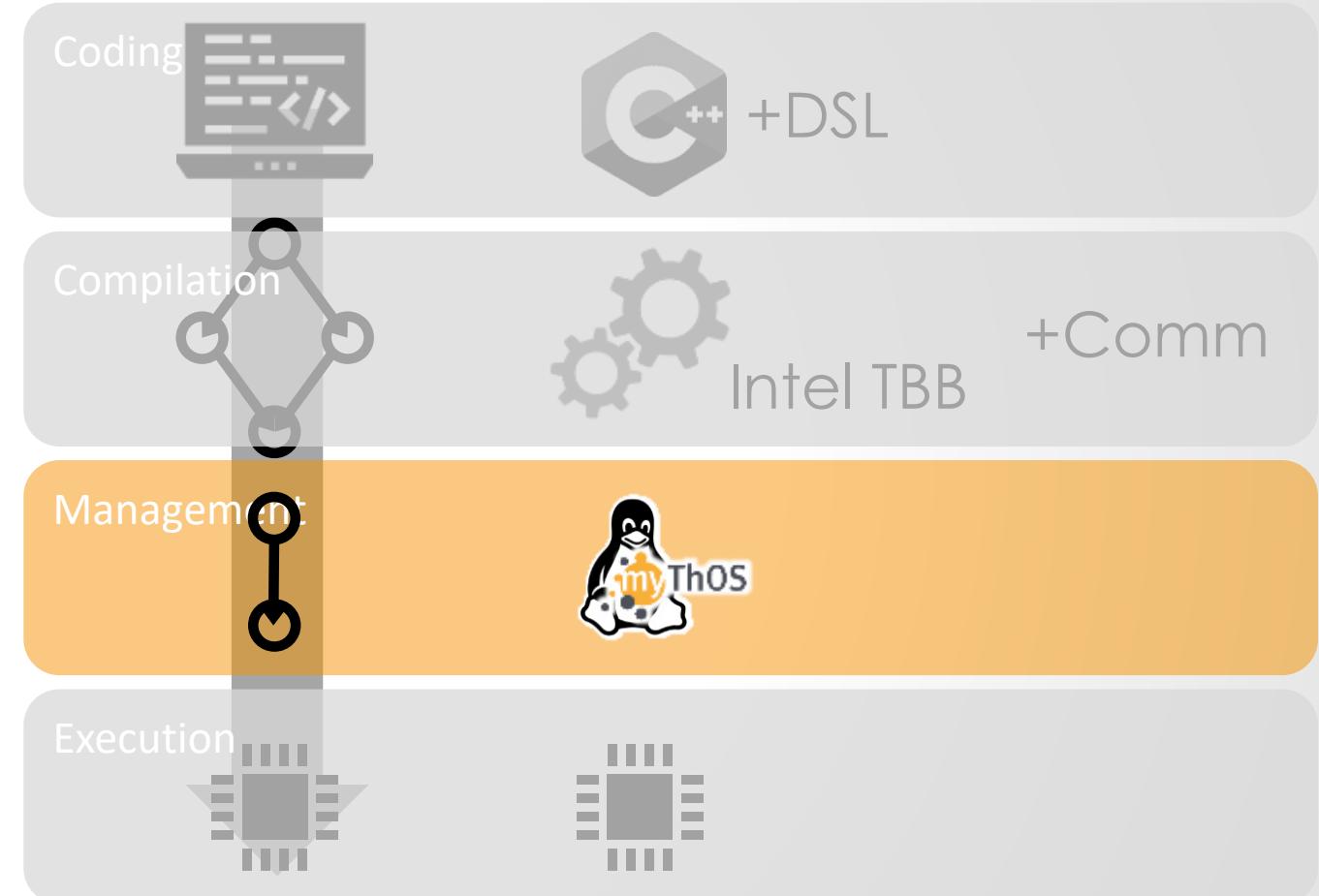
# Actual Technology & Status Node Types



# Actual Technology & Status Task Description



# Actual Technology & Status MyThOS + Linux



# Actual Technology & Status MyThOS + Linux

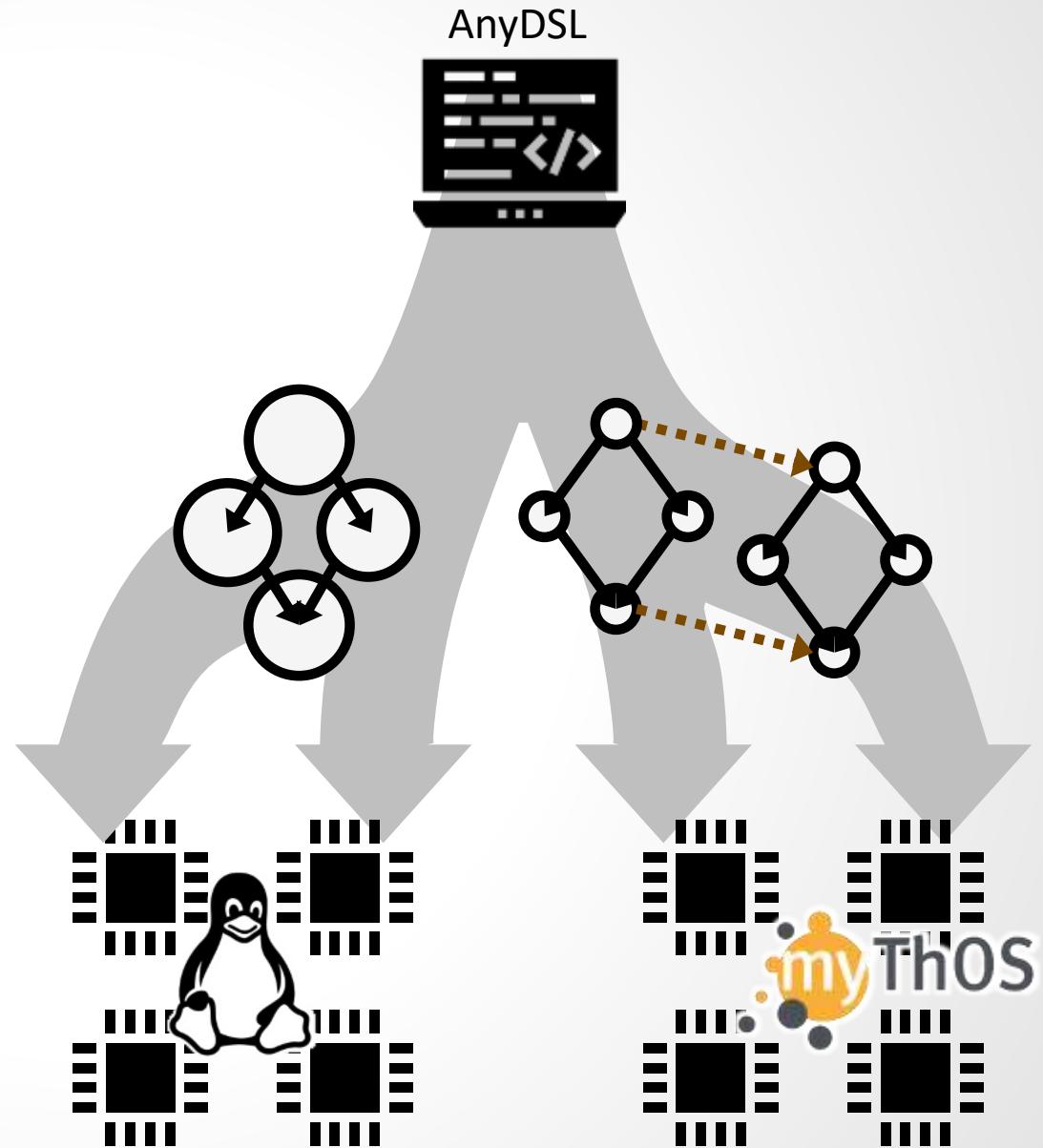


- monolithic
- creates considerable overhead
- not designed for scalability and HPC
- more difficult to adjust to new environments
- but has a network driver ;)

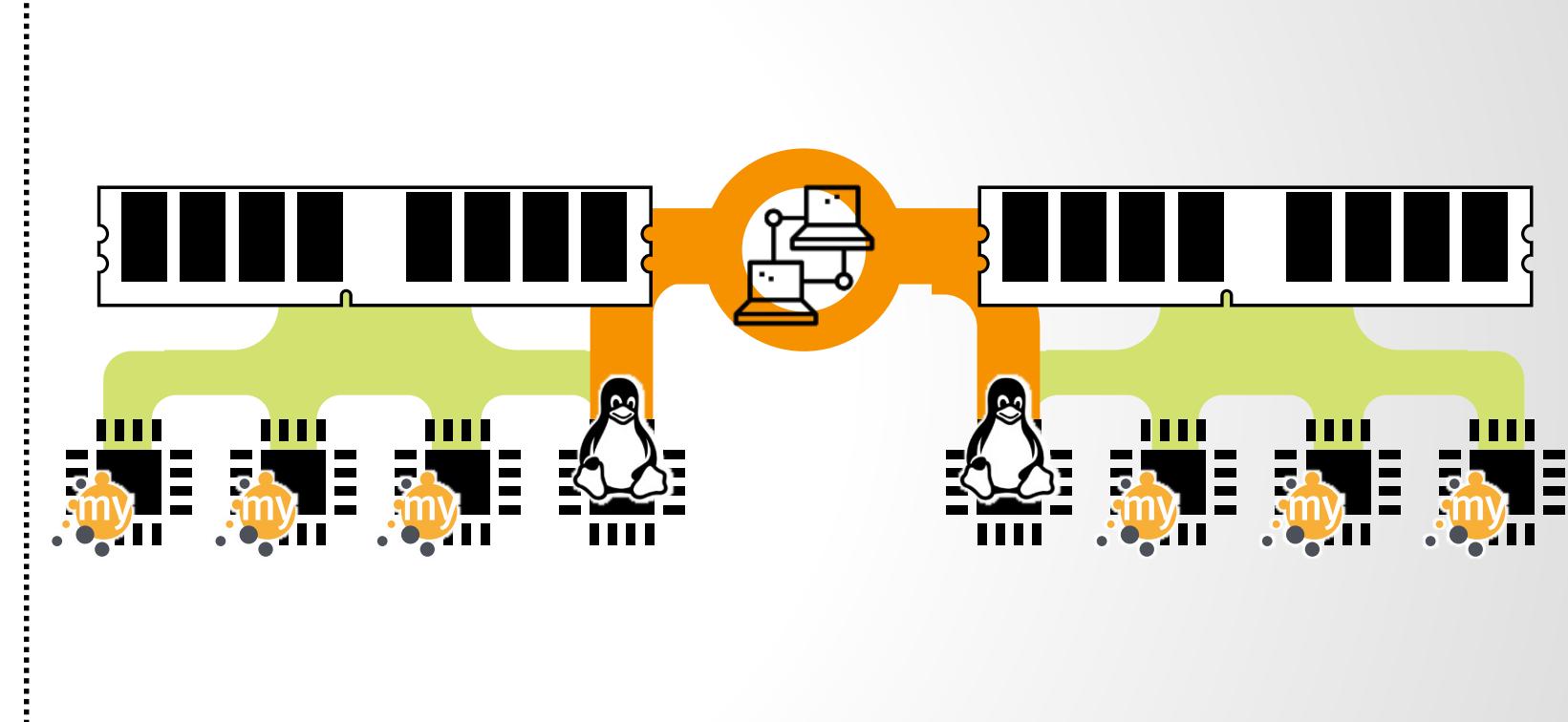


- minimal, modular OS
- provides up to 20x faster thread management than Linux => smaller tasks possible
- offers distributed memory management
- and distributed scheduling
- still missing network(!)

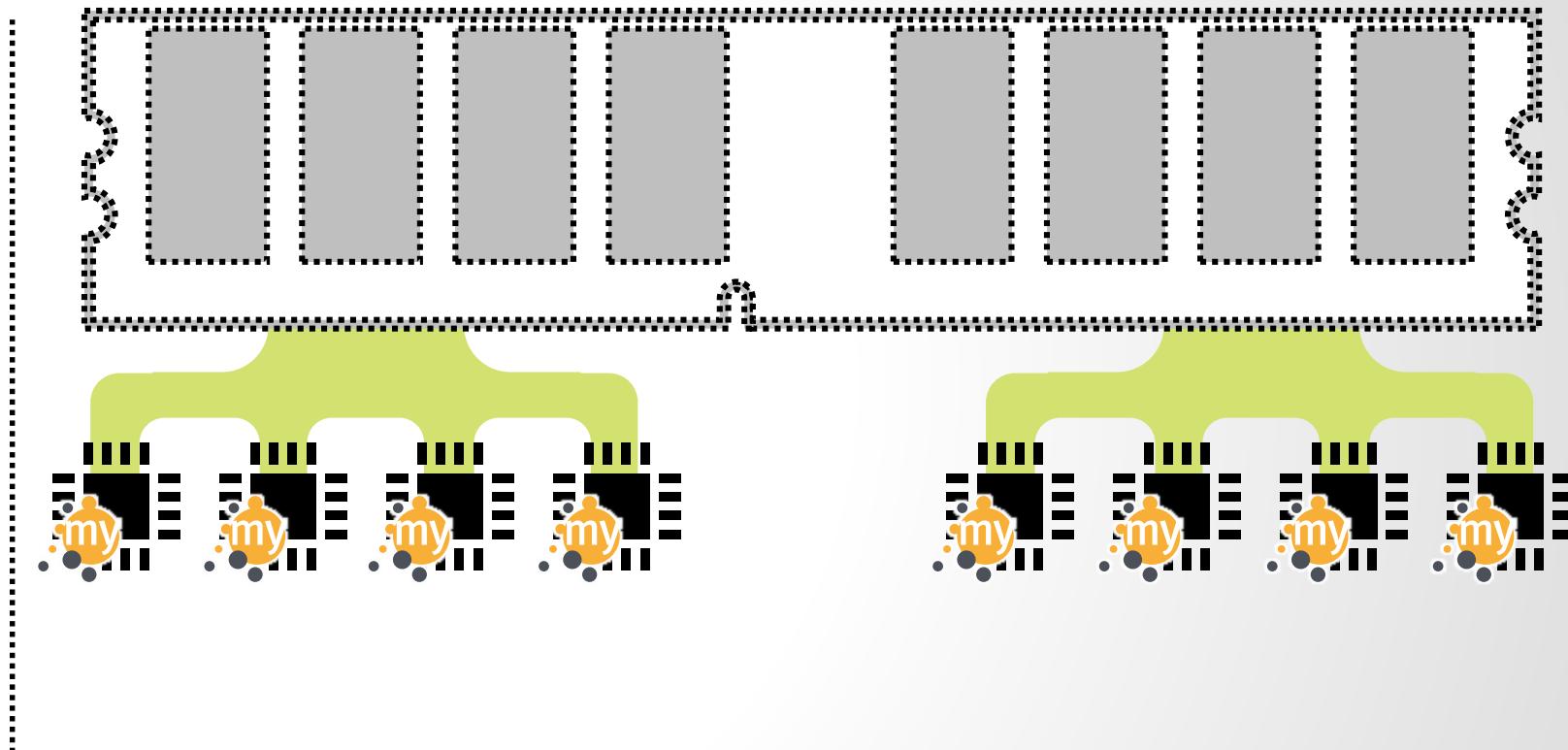
Actual  
Technology &  
Status  
MyThOS +  
Linux



# Actual Technology & Status MyThOS + Linux

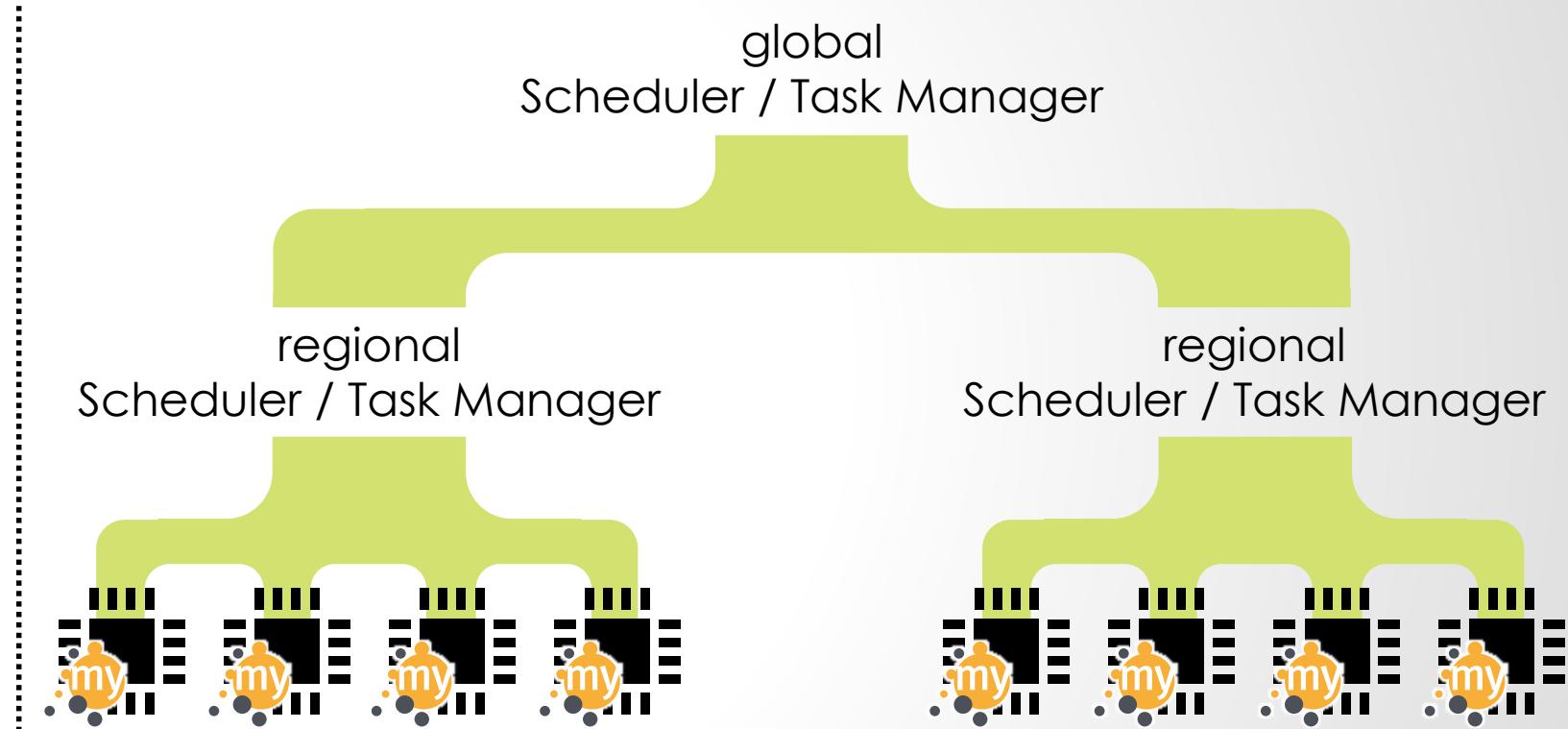


Actual Technology  
& Status  
Virtually  
Shared  
Distributed  
Memory

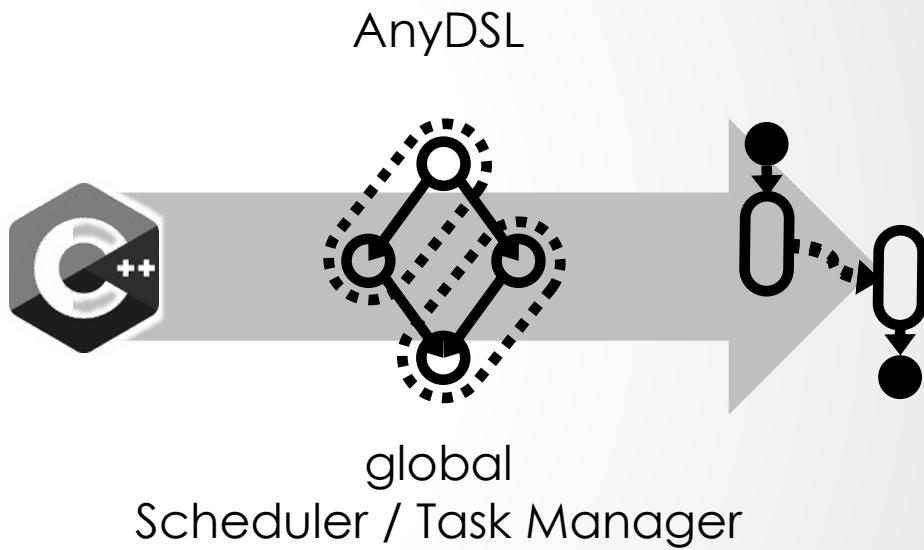


# Actual Technology & Status

## Hierarchical Scheduling



# Actual Technology & Status Flowgraph



# Actual Technology & Status Flowgraph Unrolling

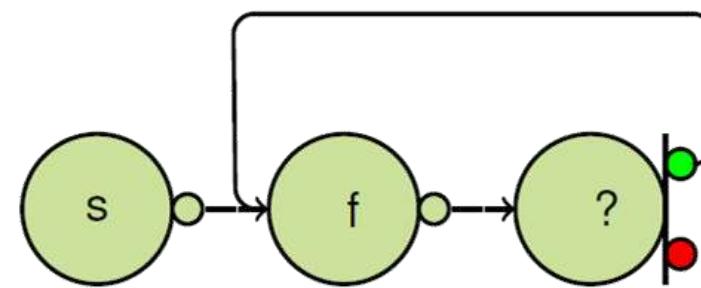


Figure: Flow Graph example

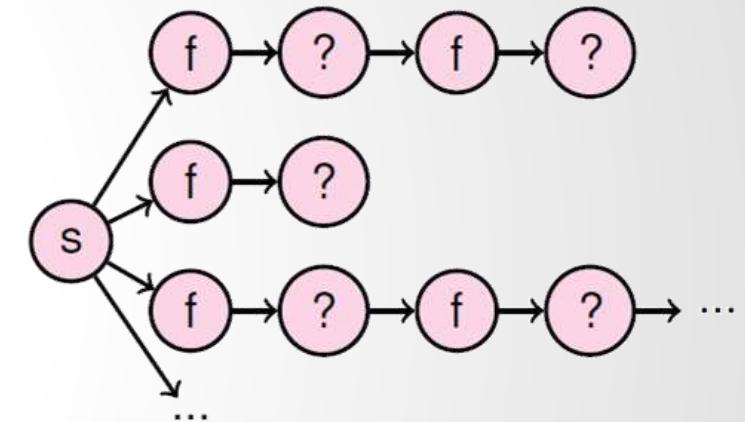
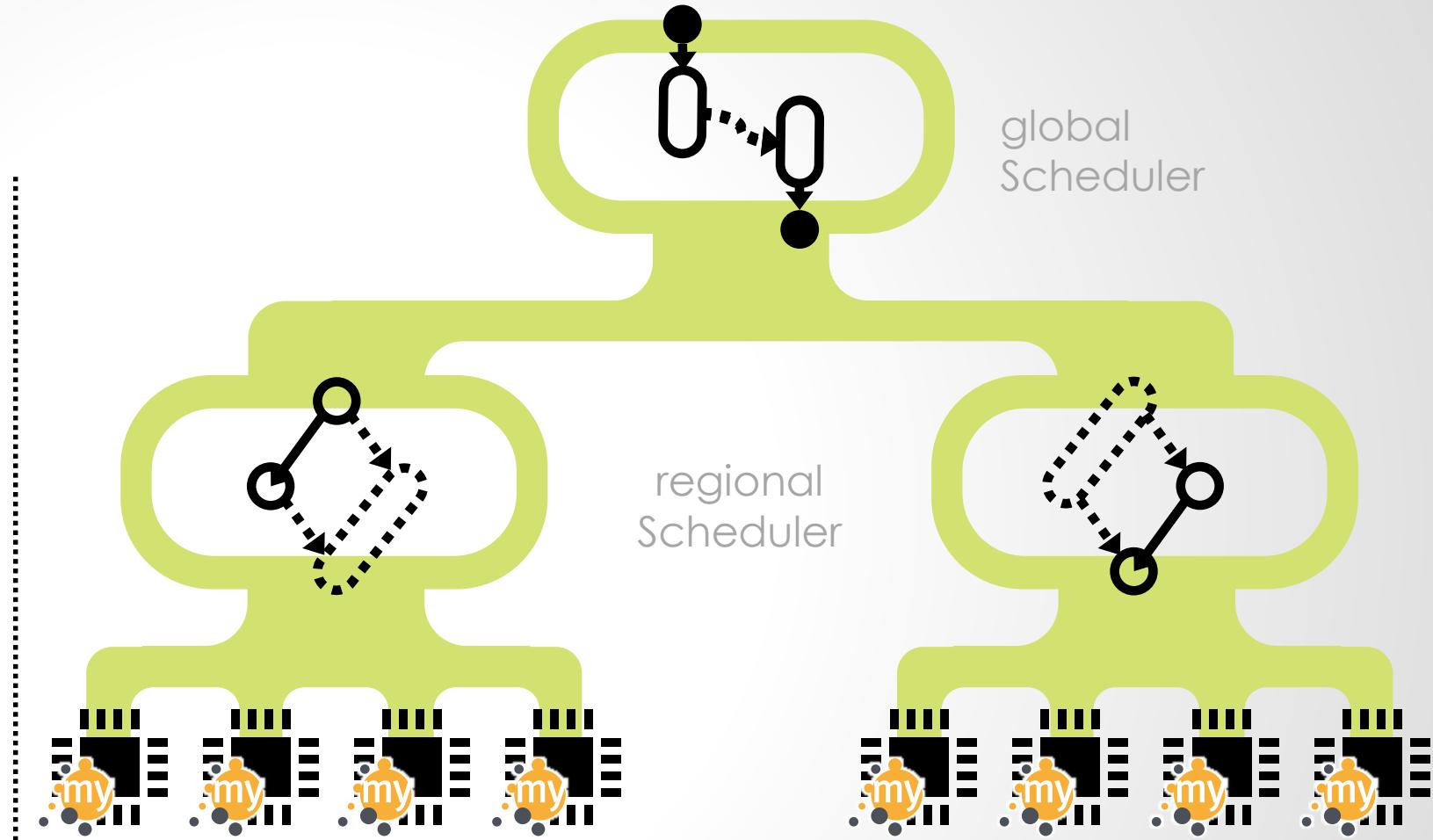
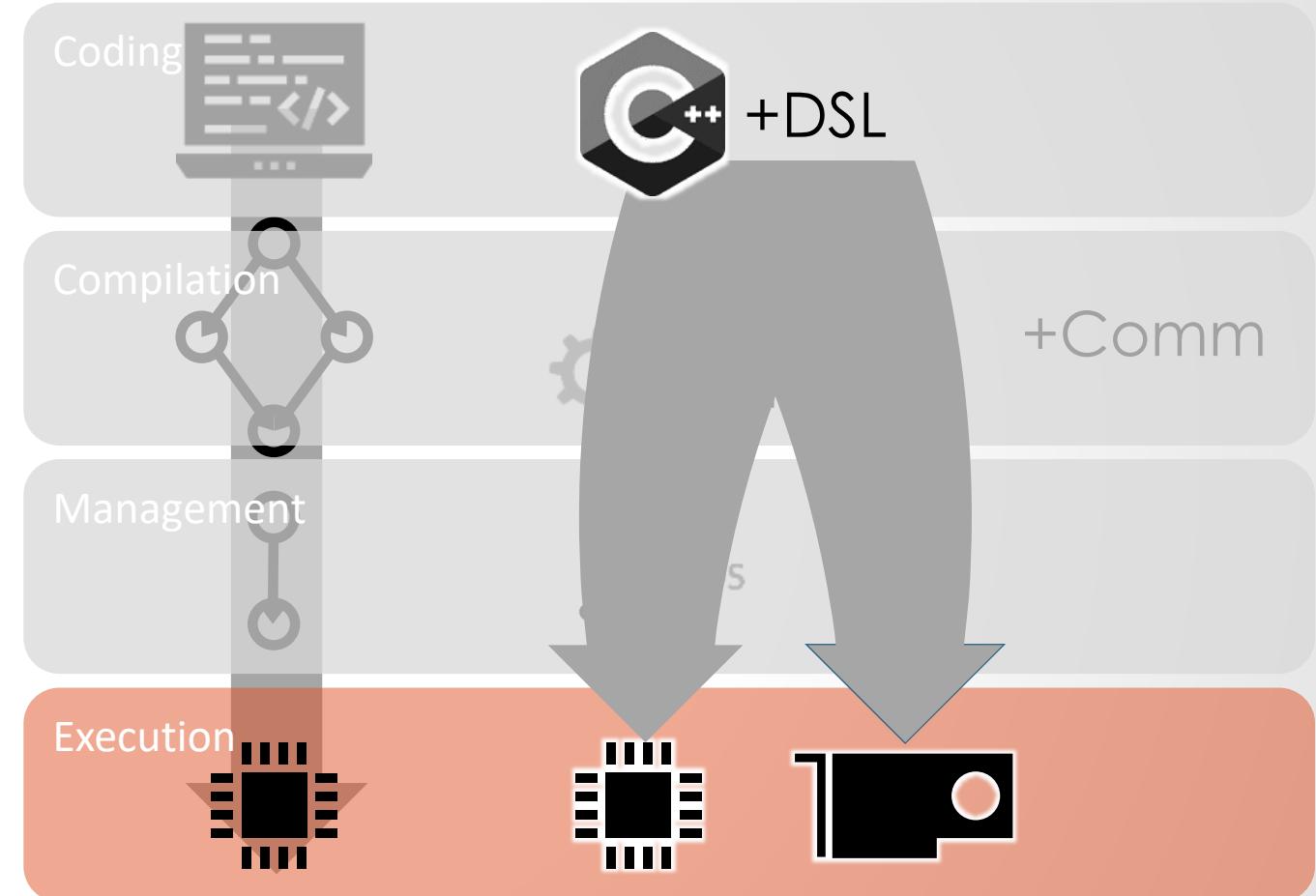


Figure: Unrolled Task Graph (DAG)

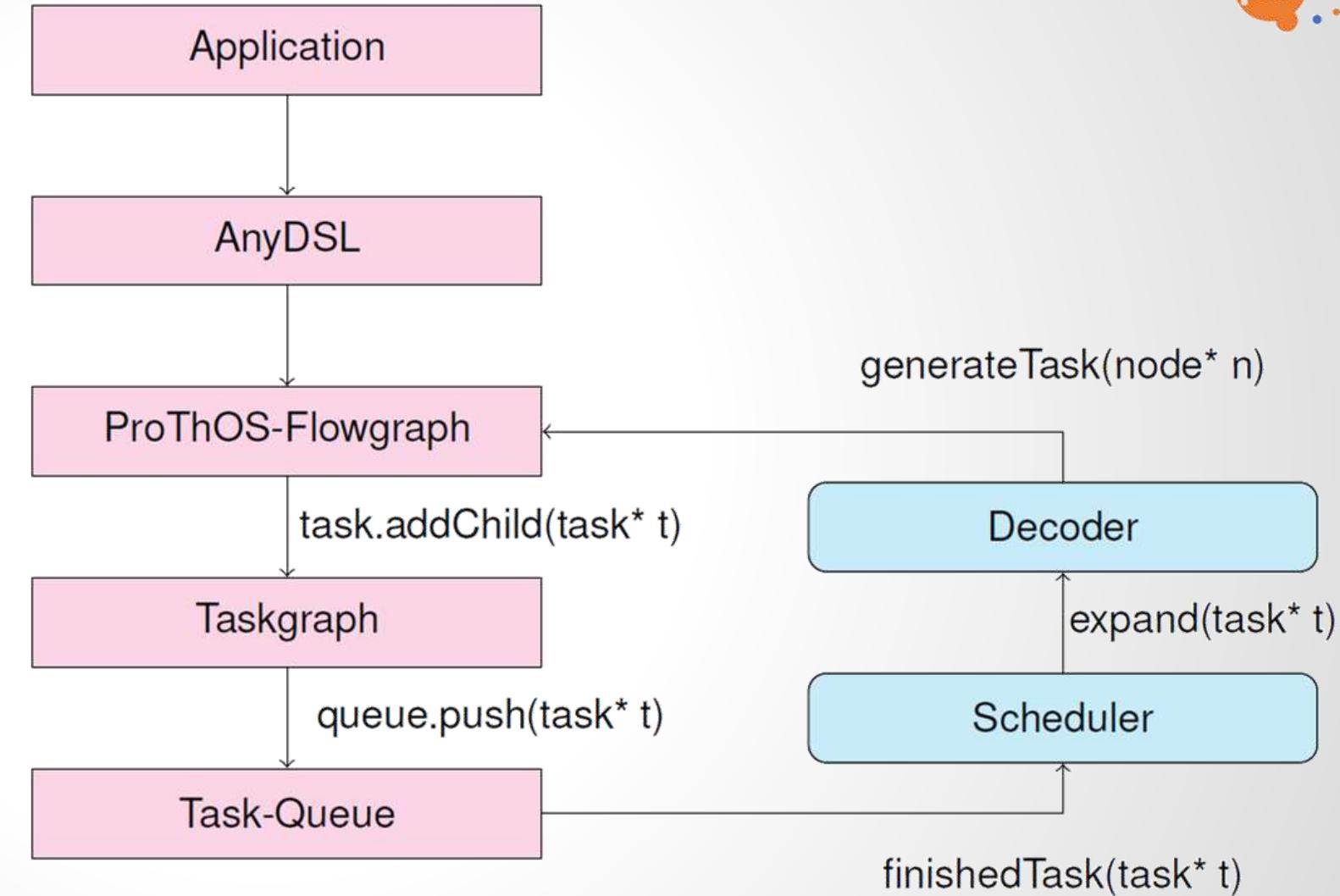
# Actual Technology & Status Flowgraph Scheduling



# Actual Technology & Status Destination Resources



# ProThOS Lifecycle Overview



# Use Case Scenario

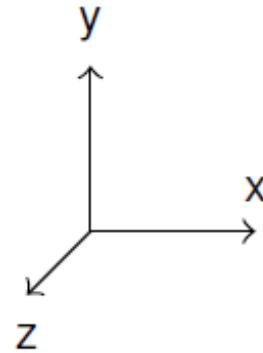
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Example

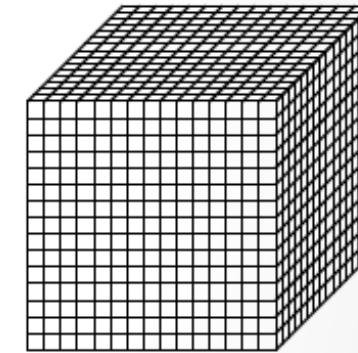
# ODTLES

## Überlagerte Gitter

- Navier-Stokes-Gleichung  $\frac{\partial \mathbf{V}}{\partial t} + \nabla(\mathbf{V} \circ \mathbf{V}) = -\frac{1}{\rho} \nabla p + \nu \nabla^2 \mathbf{V}$



DNS:  $N_{DNS}^3$

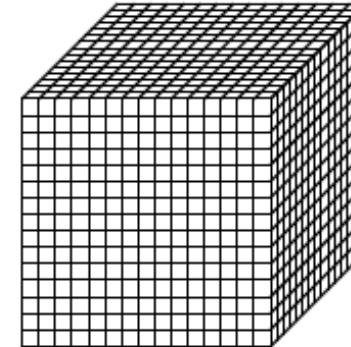


# ODTLES

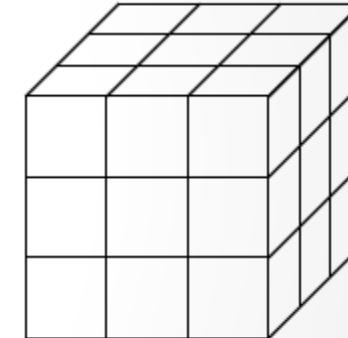
## Überlagerte Gitter

- Navier-Stokes-Gleichung  $\frac{\partial \mathbf{V}}{\partial t} + \nabla(\mathbf{V} \circ \mathbf{V}) = -\frac{1}{\rho} \nabla p + \nu \nabla^2 \mathbf{V}$
- Großskalige Advektions- und Diffusionseffekte: LES-Gitter

DNS:  $N_{DNS}^3$

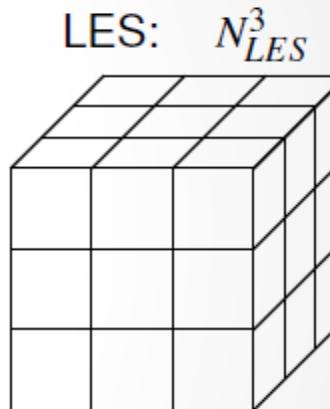
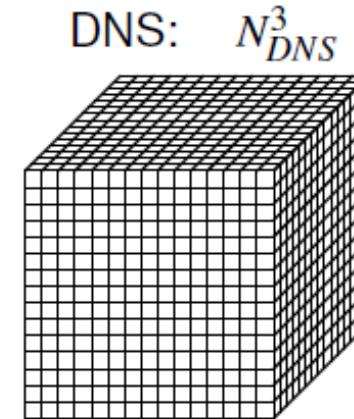


LES:  $N_{LES}^3$

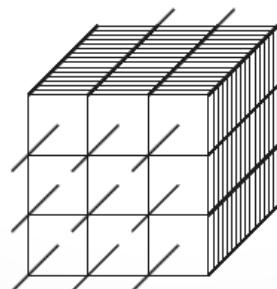


# ODTLES Überlagerte Gitter

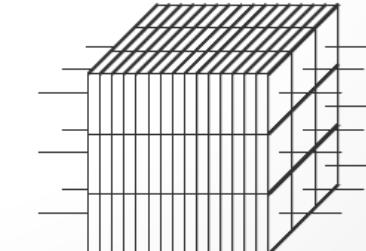
- Navier-Stokes-Gleichung  $\frac{\partial \mathbf{V}}{\partial t} + \nabla(\mathbf{V} \circ \mathbf{V}) = -\frac{1}{\rho} \nabla p + \nu \nabla^2 \mathbf{V}$
- Großskalige Advektions- und Diffusionseffekte: LES-Gitter
- Feinskalige Turbulenz: Advektion entlang 1D Linien in x,y,z-Richtung



ODTLES  $\rightarrow 3N_{LES}^2 N_{DNS}$



ODTLES-Gitter Z



ODTLES-Gitter X

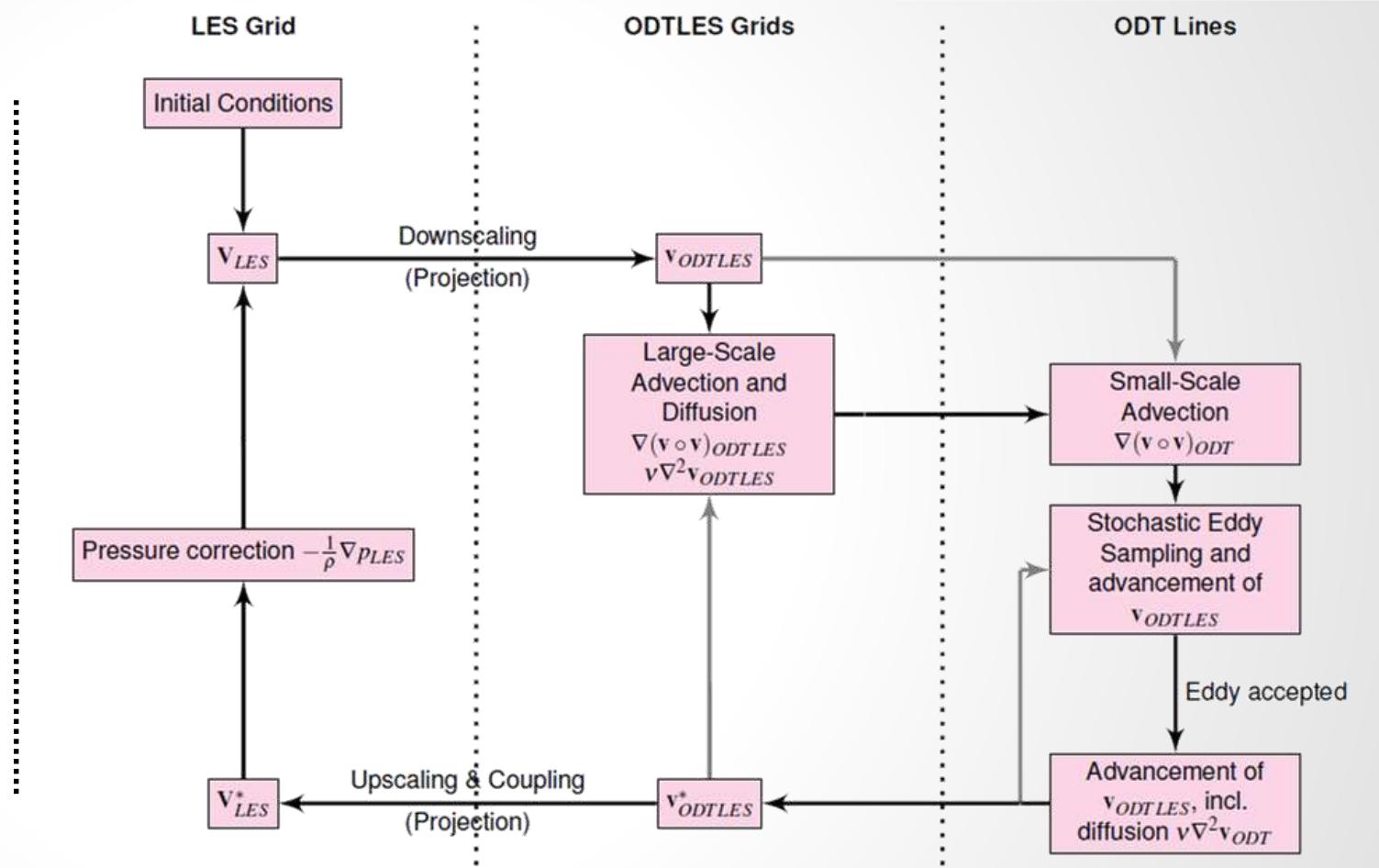


ODTLES-Gitter Y

# ODTLES Algorithmus

1. Eingabe der Anfangsbedingungen (LES)
2. Projektion der LES-Geschwindigkeiten auf die drei ODTLES-Gitter (x,y,z)
3. Bestimmung der großskaligen Advektions- und Diffusionseffekte in jedem ODTLES-Gitter
4. Stochastisches Eddy-Sampling in jeder ODT-Linie
5. Akzeptanz von Wirbeln generiert Beitrag zum zeitlichen Verlauf der ODTLES-Geschwindigkeit
6. Projektion der neuen ODTLES-Geschwindigkeiten auf das LES-Gitter
7. Druckkorrektur der LES-Geschwindigkeiten (Poisson-Problem)
8. Wiederholung von Schritt 2.-7. bis zum Ende der Simulationszeit

# ODTLES Algorithmus



# Conclusion

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Status and Next Steps

# Current Status

- ODTLES:
  - Portierung von Fortran nach C++
  - tw. AnyDSL, Verifikation
- AnyDSL:
  - JIT von C++ aus aufrufen
  - Callbacks mit Closures
  - Interface zu C++ Flow-Graphen, Stincilla Stencil-Pipelines
- Runtime:
  - Ausrollen des Flow-Graph zu azyklischem Taskgraph
- Runtime:
  - Shared-memory Prototyp auf Linux und MyThOS
- MyThOS:
  - libc++ und tw. libc Unterstützung für Anwendungen

# Next Steps

- ODTLES:
  - Beschreibung als Flow-Graph
- AnyDSL:
  - Closures mit Parametern
  - Datentransfers im Flow-Graph
- Runtime:
  - Integration zusätzlicher Knotentypen in den Flowgraph
- Scheduling:
  - Unterstützung für GPUs, Cluster-Computer



Find out more about ProThOS on  
<https://manythreads.github.io/prothos/>  
or get in touch

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